

# General Game Playing

## *Game Factoring*

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# Game Factoring

*Game factoring* (aka *game decomposition*) is the process of discovering independent subgames inside larger games.

Techniques so far provide polynomial improvement  
Factoring can provide exponential improvement

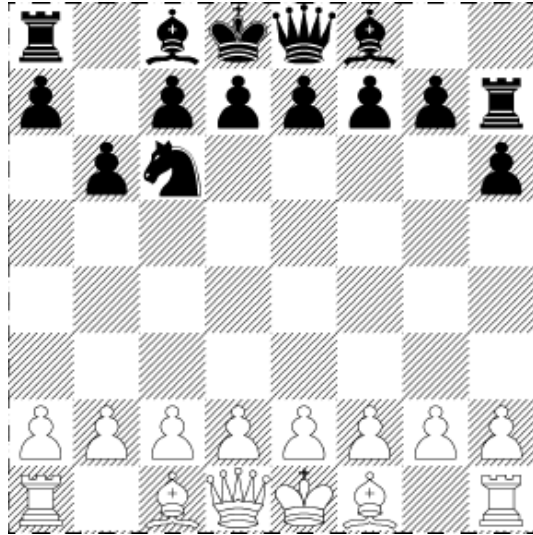
Trade-off - cost of factoring vs savings

*Sometimes* cost proportional to size of description

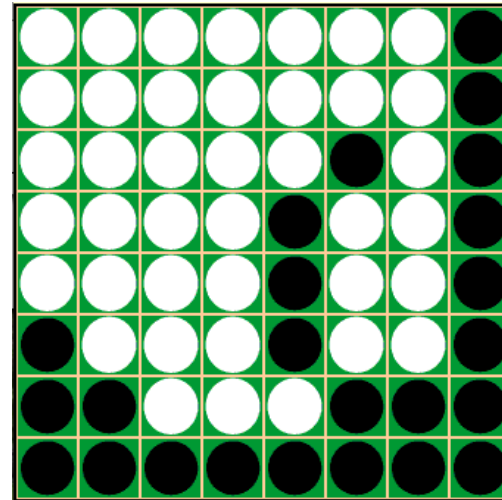
*Sometimes* savings proportional to size of game tree

# Hodgepodge

Hodgepodge = Chess + Othello



Branching factor:  $a$



Branching factor:  $b$

Analysis of joint game:

Branching factor as given to players:  $a*b$

Fringe of tree at depth  $n$  as given:  $(a*b)^n$

Fringe of tree at depth  $n$  factored:  $a^n + b^n$

# Double Tic Tac Toe

X	O	X
	O	
O		X

X		
X	O	O
O		X

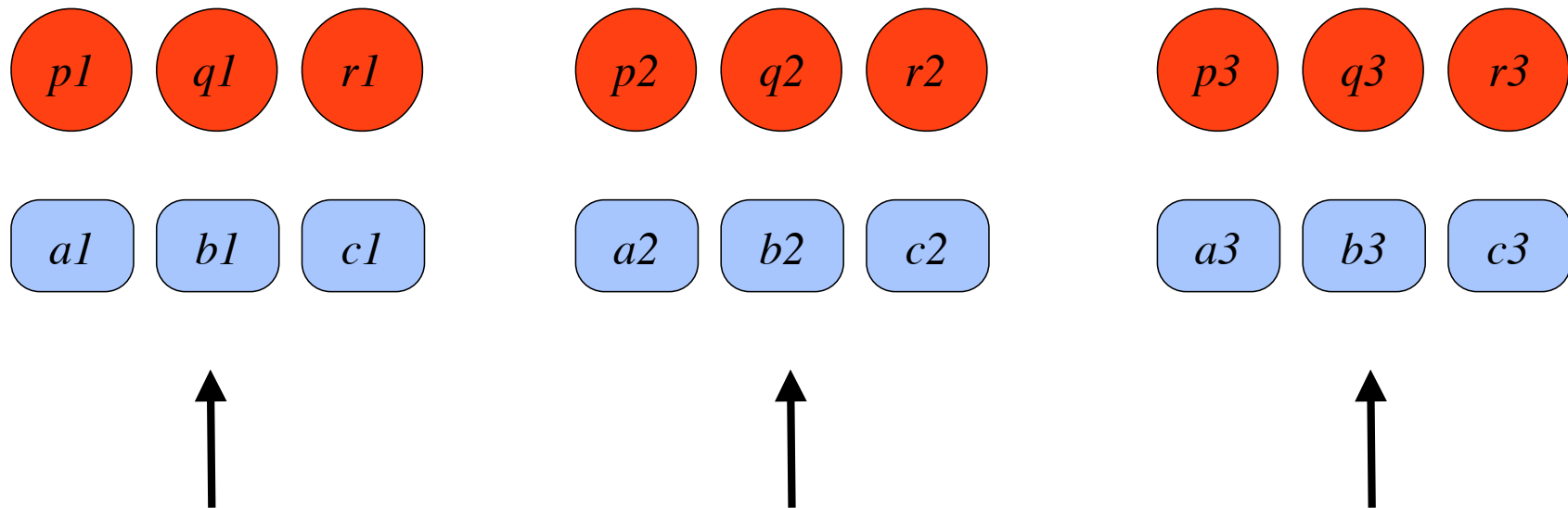
Analysis of joint game:

Branching factor: 81, 64, 49, 36, 25, 16, 9, 4, 1

Branching factor: 9, 8, 7, 6, 5, 4, 3, 2, 1

# Best Buttons and Lights

Overall terminal and goals defined as disjunctions of individual terminals and goals.



Terminating one group terminates entire game.

# Playing Factored Best Games

Technique:

- (1) Play each factor to get a move and a score.
- (2) Select subgame/move with best score.

Cost:

$$\text{cost}(\text{game}_1) + \dots + \text{cost}(\text{game}_k) \ll \text{cost}(\text{game})$$

Nodes searched by Minimax on 9 board BB&L

*with* factoring: 3276

*without* factoring: 14348907

# Grounding

# Grounding

```
next(p(X)) :- does(robot,a(X)) & ~true(p(X))  
next(p(X)) :- does(robot,b(X)) & true(q(X))  
next(p(X)) :- true(p(X)) & ~does(robot,a(X)) & ~does(robot,b(X))
```



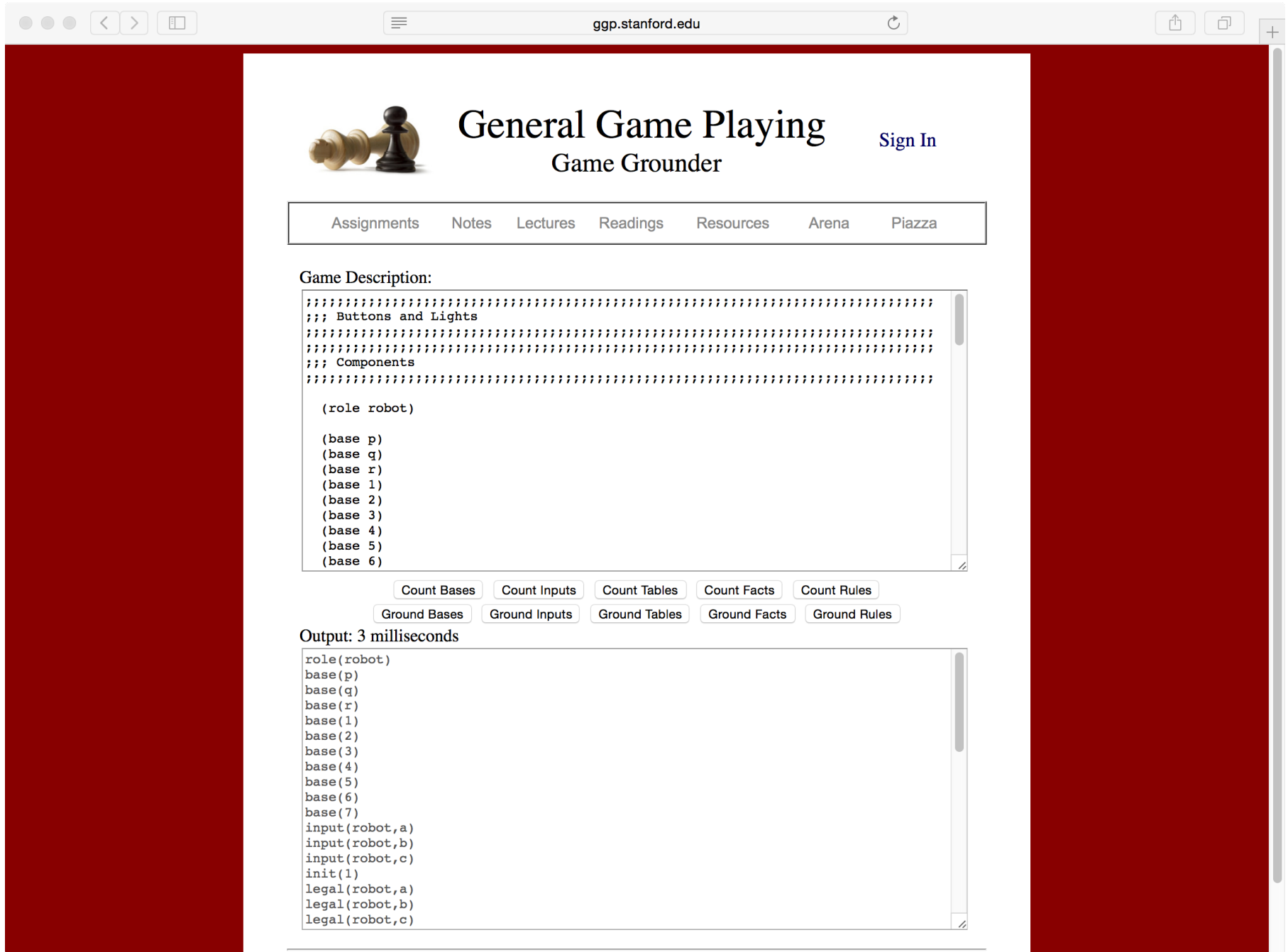
```
next(p(1)) :- does(robot,a(1)) & ~true(p(1))  
next(p(1)) :- does(robot,b(1)) & true(q(1))  
next(p(1)) :- true(p(1)) & ~does(robot,a(1)) & ~does(robot,b(1))
```

```
next(p(2)) :- does(robot,a(2)) & ~true(p(2))  
next(p(2)) :- does(robot,b(2)) & true(q(2))  
next(p(2)) :- true(p(2)) & ~does(robot,a(2)) & ~does(robot,b(2))
```

```
next(p(3)) :- does(robot,a(3)) & ~true(p(3))  
next(p(3)) :- does(robot,b(3)) & true(q(3))  
next(p(3)) :- true(p(3)) & ~does(robot,a(3)) & ~does(robot,b(3))
```



# Grounding Tool



The screenshot shows a web browser window with the URL `ggp.stanford.edu`. The page features a navigation menu with items: Assignments, Notes, Lectures, Readings, Resources, Arena, and Piazza. The main content area is titled "General Game Playing Game Grounder" and includes a "Sign In" link. Below the title is a "Game Description:" section containing a text area with Prolog code. The code defines a robot role and six base positions. Below the code are buttons for "Count Bases", "Count Inputs", "Count Tables", "Count Facts", and "Count Rules". Underneath these are buttons for "Ground Bases", "Ground Inputs", "Ground Tables", "Ground Facts", and "Ground Rules". The "Output: 3 milliseconds" section shows the resulting grounded Prolog code, which includes the original definitions plus additional facts for the robot's input and legal actions.

General Game Playing  
Game Grounder [Sign In](#)

Assignments Notes Lectures Readings Resources Arena Piazza

Game Description:

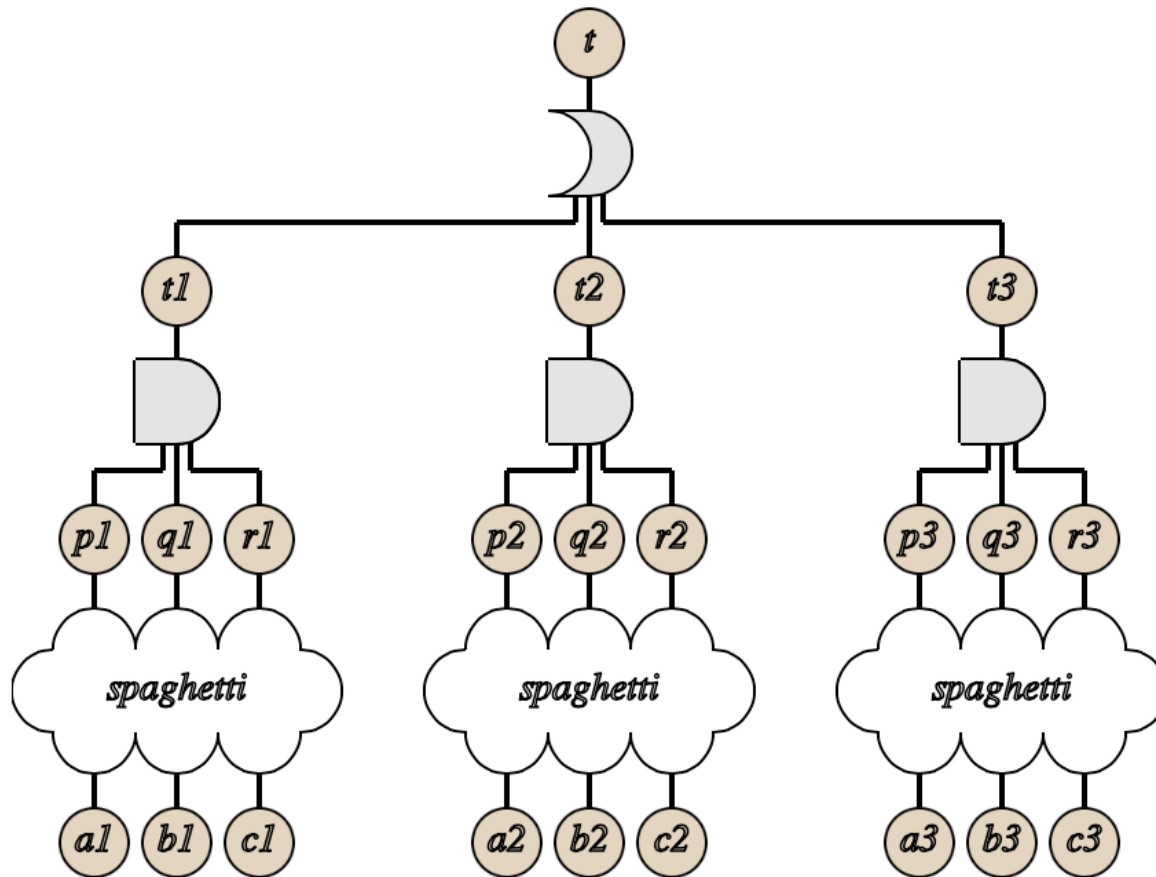
```
;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;  
;;; Buttons and Lights  
;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;  
;;; Components  
;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;;  
  
(role robot)  
  
(base p)  
(base q)  
(base r)  
(base 1)  
(base 2)  
(base 3)  
(base 4)  
(base 5)  
(base 6)
```

Count Bases Count Inputs Count Tables Count Facts Count Rules  
Ground Bases Ground Inputs Ground Tables Ground Facts Ground Rules

Output: 3 milliseconds

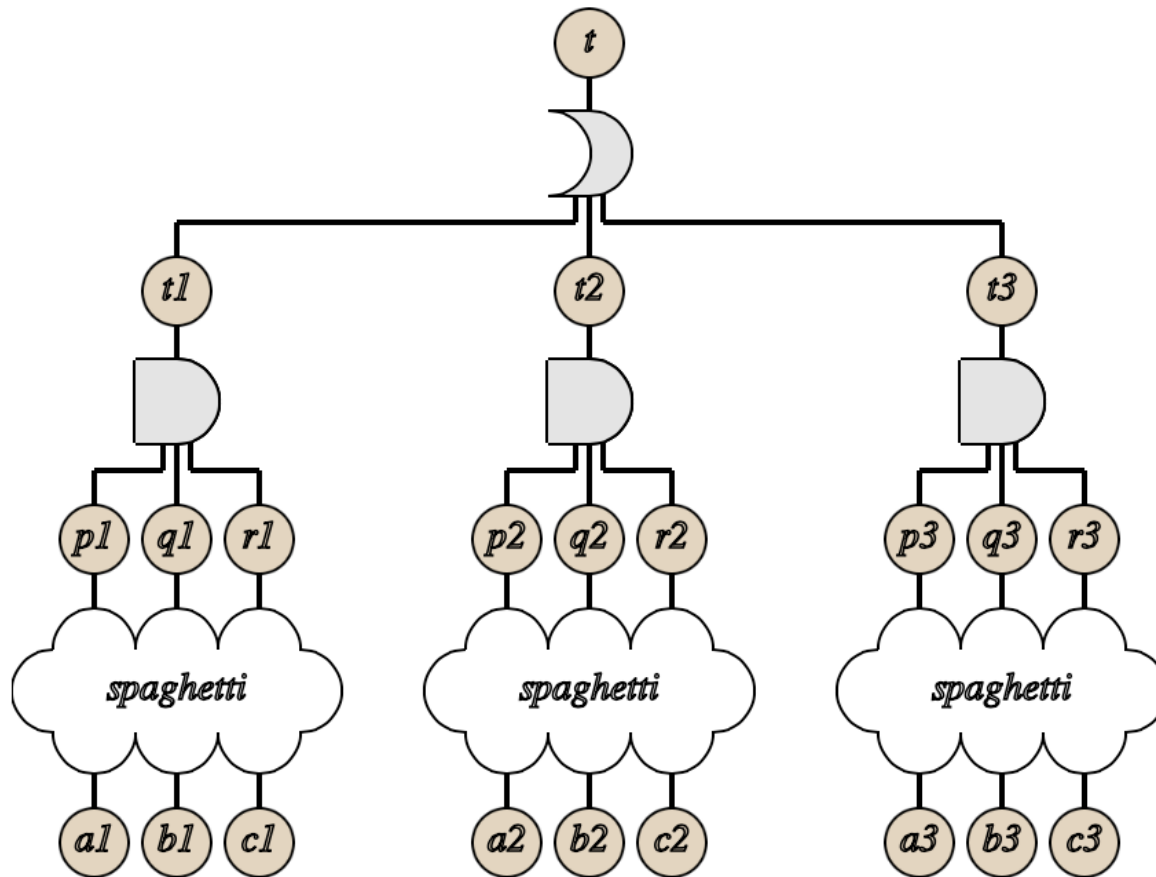
```
role(robot)  
base(p)  
base(q)  
base(r)  
base(1)  
base(2)  
base(3)  
base(4)  
base(5)  
base(6)  
base(7)  
input(robot,a)  
input(robot,b)  
input(robot,c)  
init(1)  
legal(robot,a)  
legal(robot,b)  
legal(robot,c)
```

# Propnet for Best Buttons and Lights



# Factoring

# Propnet for Best Buttons and Lights



# Base and Input Propositions

## Bases:

```
base(p(X)) :- index(X)
base(q(X)) :- index(X)
base(r(X)) :- index(X)
index(1)
index(2)
index(3)
```

## Inputs:

```
input(robot,a(X)) :- index(X)
input(robot,b(X)) :- index(X)
input(robot,c(X)) :- index(X)
```

## Results:

```
true(p(1))
true(p(2))
true(p(3))
true(q(1))
true(q(2))
true(q(3))
true(r(1))
true(r(2))
true(r(3))
```

```
does(robot(a(1)))
does(robot(a(2)))
does(robot(a(3)))
does(robot(b(1)))
does(robot(b(2)))
does(robot(b(3)))
does(robot(c(1)))
does(robot(c(2)))
does(robot(c(3)))
```

# Compute Residues

Compute all  $\text{true}(p)$  and  $\text{does}(role, a)$  factoids needed to compute  $\text{next}(q)$  and  $\text{legal}(role, a)$  for each base proposition  $q$  and each input  $a$ . These sets are called *residues*.

Rules for  $\text{next}(p(1))$ :

$\text{next}(p(1)) \text{ :- does(robot, a(1)) \& \sim true(p(1))}$

$\text{next}(p(1)) \text{ :- does(robot, b(1)) \& true(q(1))}$

$\text{next}(p(1)) \text{ :- true(p(1)) \& \sim does(robot, a(1)) \& \sim does(robot, b(1))}$

Residue for  $\text{next}(p(1))$ :

$\{\text{true}(p(1)), \text{true}(q(1), \text{does}(\text{robot}, a(1))), \text{does}(\text{robot}, b(1))\}$

# Merge Overlapping Residues

```
{true(p(1)), true(q(1),  
  does(robot,a(1))), does(robot,b(1))}  
{true(q(1)), true(p(1)), true(r(1)),  
  does(robot,b(1))), does(robot,c(1))}  
{true(r(1)), true(q(1)), does(robot,c(1))}
```

```
{true(p(2)), true(q(2),  
  does(robot,a(2))), does(robot,b(2))}  
{true(q(2)), true(p(2)), true(r(2)),  
  does(robot,b(2))), does(robot,c(2))}  
{true(r(2)), true(q(2)), does(robot,c(2))}
```

```
{step(1), step(2), ..., step(7)}
```

```
{true(p(1)), true(q(1), true(r(1))  
  does(robot,a(1))), does(robot,b(1), does(robot,c(1)))}  
{true(p(2)), true(q(2), true(r(2))  
  does(robot,a(2))), does(robot,b(2), does(robot,c(2)))}  
{step(1), step(2), ..., step(7)}
```

# Disjoint Action Sets

`{does(robot,a(1)), does(robot,b(1), does(robot,c(1)))}`

`{does(robot,a(2)), does(robot,b(2), does(robot,c(2)))}`

`{does(robot,a(3)), does(robot,b(3), does(robot,c(3)))}`



# Rules Defining Legality

```
legal(robot,a(1)) :- index(1)
legal(robot,b(1)) :- index(1)
legal(robot,c(1)) :- index(1)
legal(robot,a(2)) :- index(2)
legal(robot,b(2)) :- index(2)
legal(robot,c(2)) :- index(2)
legal(robot,a(3)) :- index(3)
legal(robot,b(3)) :- index(3)
legal(robot,c(3)) :- index(3)
```

# Partition into Different Subgames

```
legal(robot,a(1)) :- index(1)  
legal(robot,b(1)) :- index(1)  
legal(robot,c(1)) :- index(1)
```

```
legal(robot,a(2)) :- index(2)  
legal(robot,b(2)) :- index(2)  
legal(robot,c(2)) :- index(2)
```

```
legal(robot,a(3)) :- index(3)  
legal(robot,b(3)) :- index(3)  
legal(robot,c(3)) :- index(3)
```

# Playing Factored Best Games

Technique:

- (1) Play each factor to get a move and a score.
- (2) Select subgame/move with best score.

Cost:

$$\text{cost}(\text{game}_1) + \dots + \text{cost}(\text{game}_k) \ll \text{cost}(\text{game})$$

Nodes searched by Minimax on 9 board BB&L

*with* factoring: 3276

*without* factoring: 14348907

# Nine Board Best Buttons and Lights

Depth	Unfactored	Factored	Time (msec)
1	28	36	14
2	757	117	45
3	20,440	360	140
4	551,881	1089	420
5	14,369,347	3276	1300
6		3437	1350

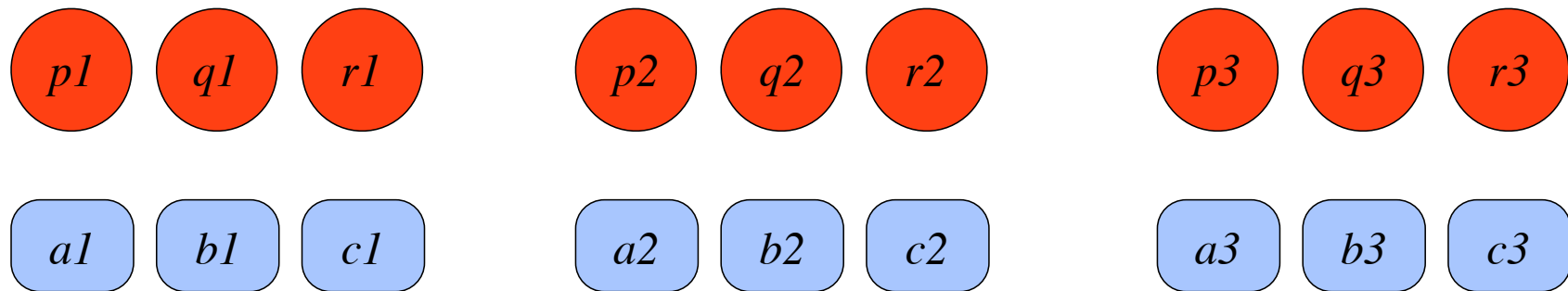
Unfactored:  $1 + 27 + 27^2 + 27^3 + 27^4 + 27^5$

Factored:  $(1 + 3 + 3^2 + 3^3 + 3^4 + 3^5) \times 9$

# Conditions

# Inertiality

Overall terminal and goals defined as disjunctions of individual terminals and goals.



Suppose one group changes on action in another group.

# Inertiality

A subgame is *inertial* iff the state of the subgame does not change when no action in the subgame is performed.

# Technique

Inertiality is computed by checking  $\text{next}(p)$  for every base proposition  $p$ .

Reduce all rules for  $\text{next}(p)$  to rules defining  $\text{next}(p)$  in terms of  $\text{true}(p)$  and  $\text{does}(r, p)$ , resolve with each other, and filter out subsumed rules.

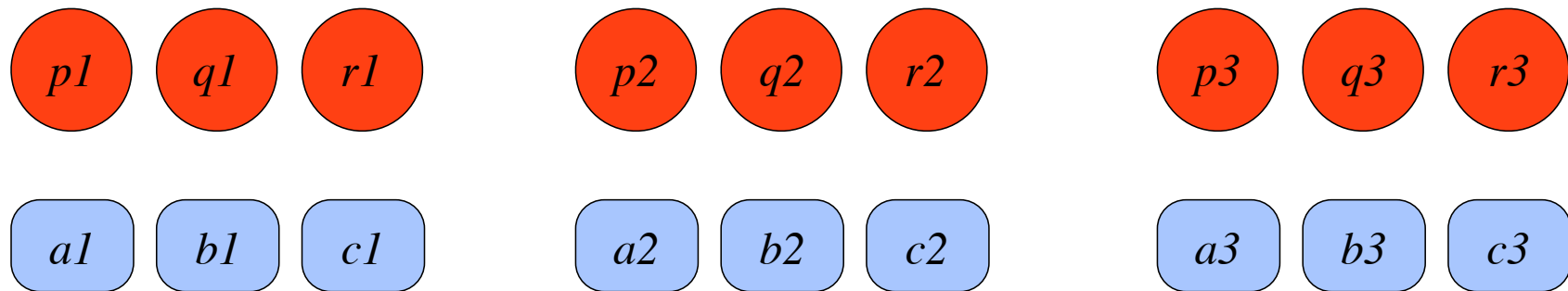
There must be a rule  $\text{next}(p) \text{ :- true}(p) \ \& \ \text{junk}$  where  $\text{junk}$  does not depend on base propositions.

There must not be any rule of the form  $\text{next}(p) \text{ :- junk}$  where  $\text{junk}$  does not include  $\text{true}(p)$ .



# No Termination

Overall terminal and goals defined as disjunctions of individual terminals and goals.



Is it possible for a group to have no terminal state?

# Termination

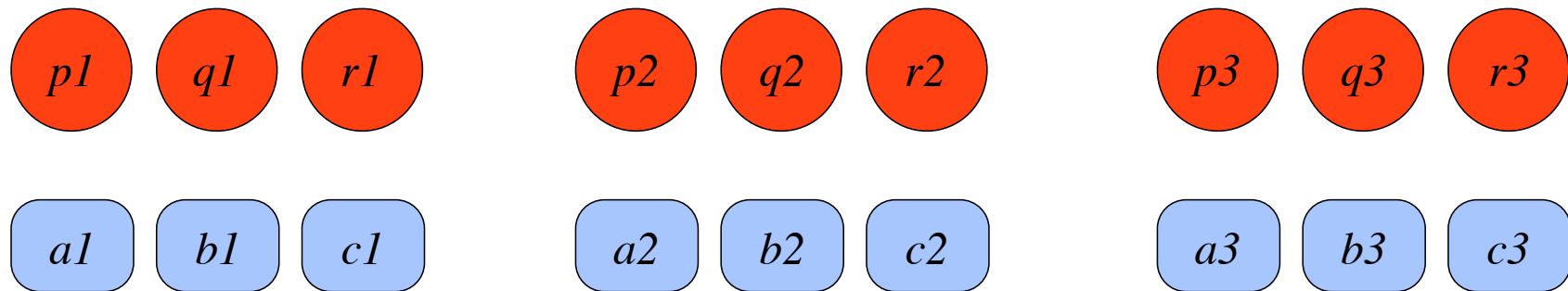
Possible for a subgame not to have a termination condition at all?

Not possible. All games must terminate. If game is inertial and some subgame does not have a termination condition, then it would be possible to play repeatedly in the game forever, contradicting the requirement for termination of all games.

Upshot: No check necessary.

# Goal Coverage

Overall terminal and goals defined as disjunctions of individual terminals and goals.



Is it possible for one group to terminate with low goal while another is non-terminal with higher goal?

# Goal Coverage

In general, it is possible for a subgame to have a high goal and no termination while another has termination and low goal.

Problem resolved if all non-zero goal states are also terminal states.

Can be checked in a manner similar to checking for inertiality.

# Validation Problem

Check that the game satisfies all of these conditions.

Almost as difficult or even more difficult than factoring.

# Brian's Method

(1) Factor

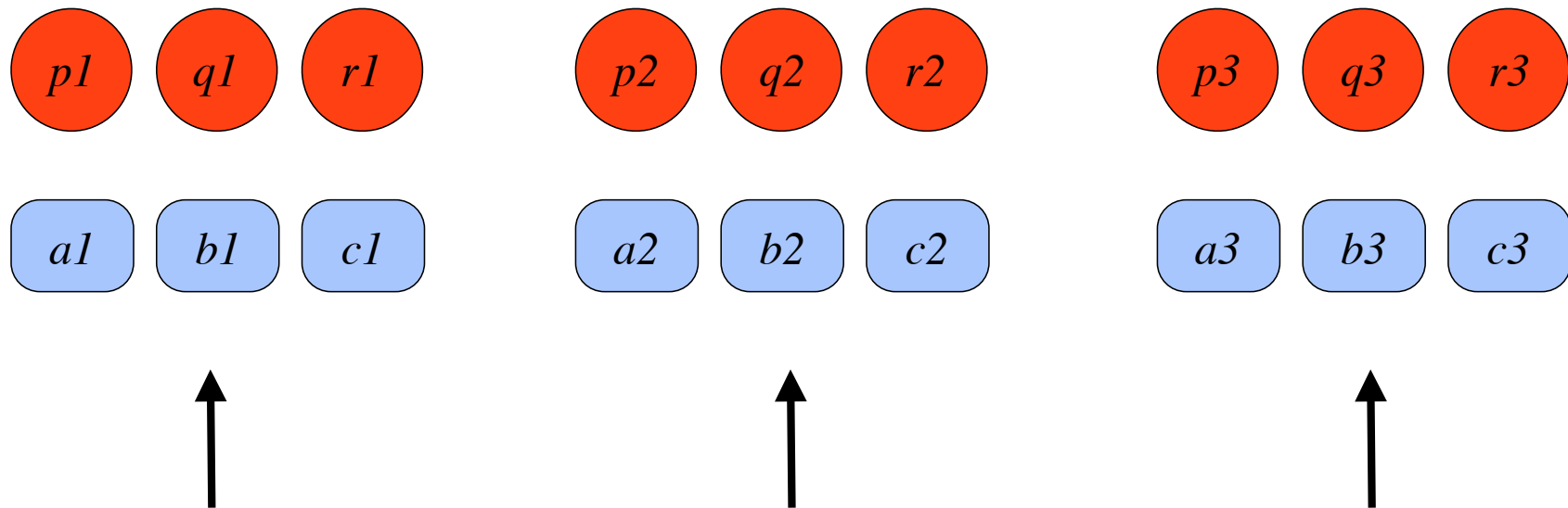
(2) Get a solution

(3) See if it works in the unfactored game.

# Other Types of Game Analysis

# Best Buttons and Lights

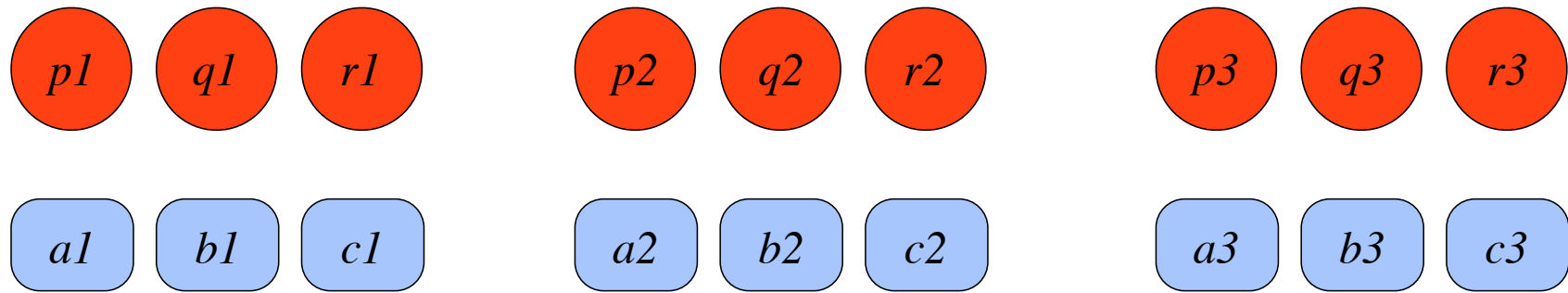
Overall terminal and goals defined as disjunctions of individual terminals and goals.



Terminating one group terminates entire game.

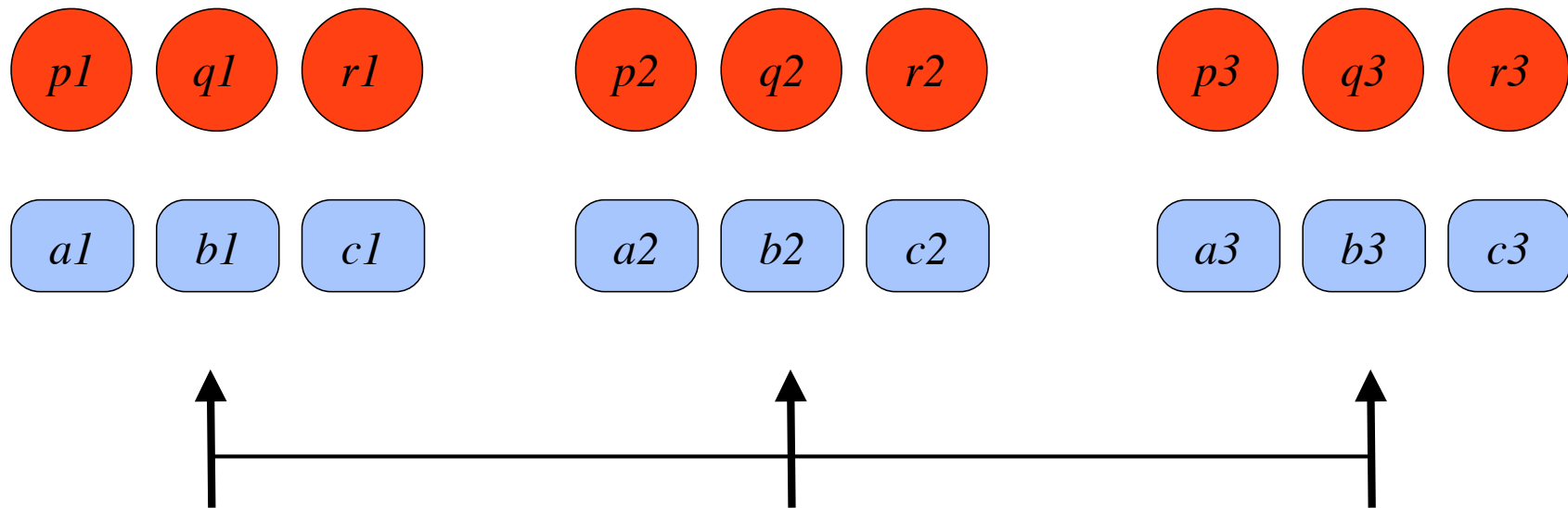


# Multiple Buttons and Lights



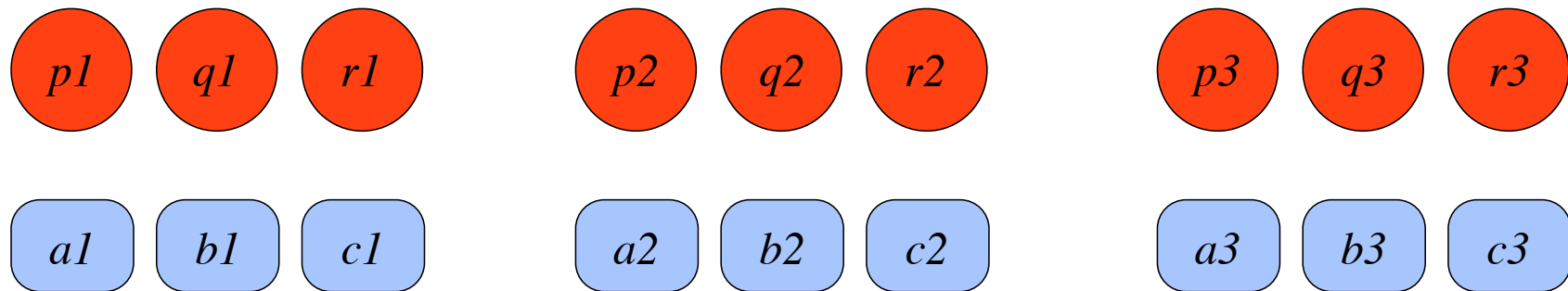
↑  
Only this group matters

# Joint Buttons and Lights



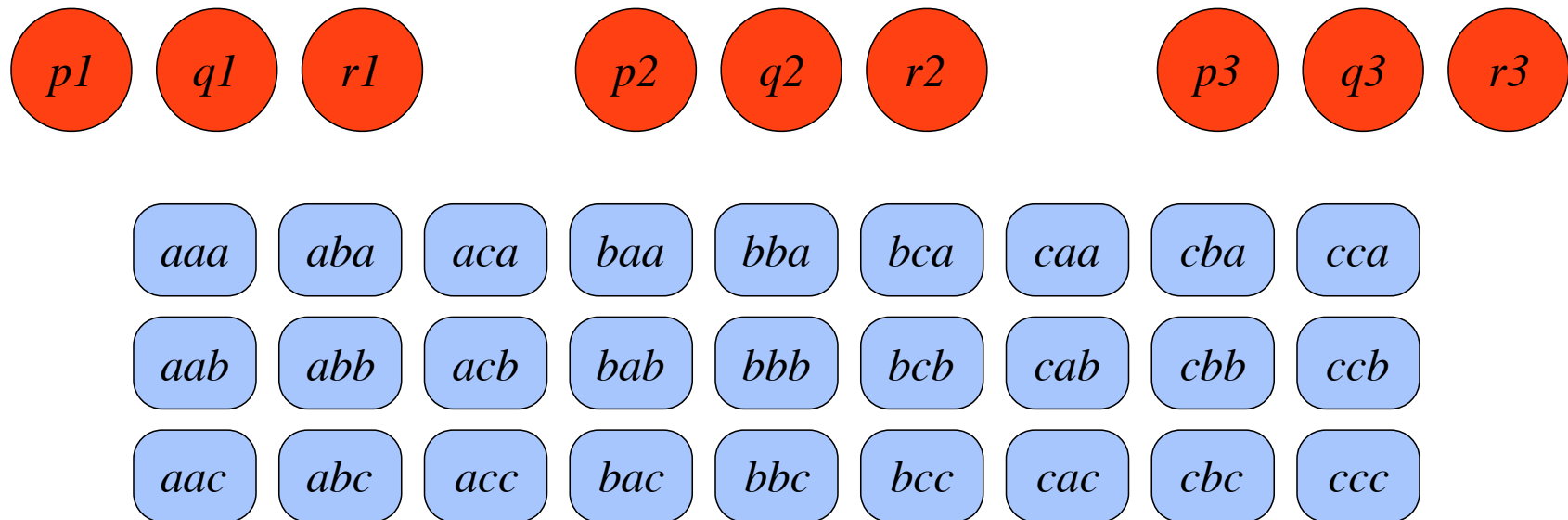
Terminates if and only if *all* groups terminate.

# Parallel Buttons and Lights

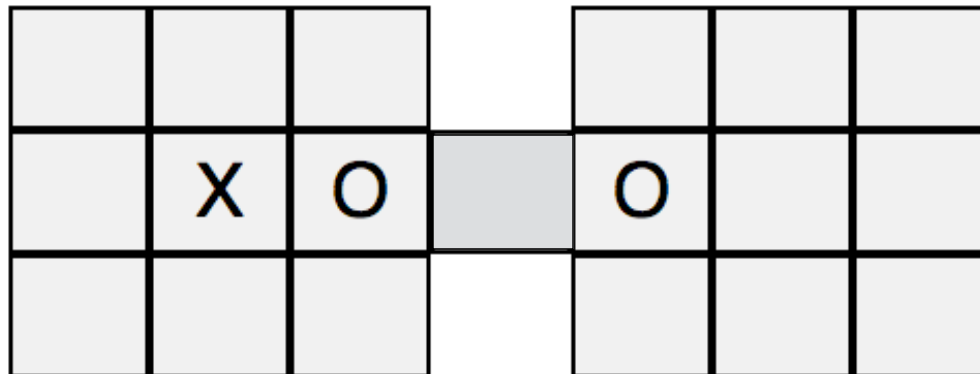


Terminal if and only if *all* groups terminal.  
Each group acted upon by different player.

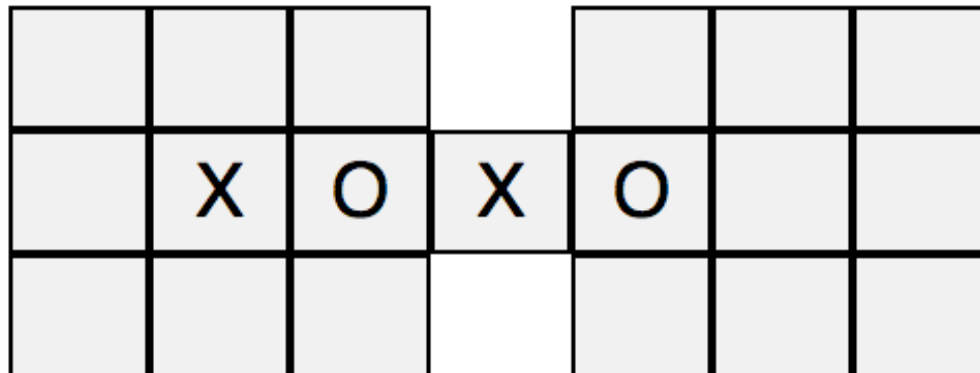
# Simultaneous Buttons and Lights



# Conditional Factoring



# Conditional Factoring



# Related Types of Game Analysis

## Bottlenecks

Series of games

each of which must terminate before next begins

## Dead State Elimination

Find states that cannot lead to acceptable outcomes

Prune whole subtrees

## Goal Monotonicity

Detect monotonicity in states

e.g. higher goal value in non-terminal states

correlated with progress toward goal



**GENERAL  
GAME  
PLAYING**





